IPASIR-UP: User Propagators for CDCL

SAT Conference, July 7, 2023 Alghero, Italy

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Usual Use of SAT Solvers



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Model enumeration, model checking, CEGAR, symbolic execution, ...

Use of Incremental SAT Solvers via User Propagators



Combinatorial problems, SMT, symmetry breaking, model enumeration, MaxSAT, ...

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IPASIR-UP: A New Interface for Interactive CDCL

Interface to support standardized interactions with the solver during solving

- Needs to be implemented in SAT solvers (only once)
- + Easy to implement and use
- + Solver independent application development
- + No more black-box SAT solving \rightarrow new potentials
- + Standardized and clean interactions

IPASIR Model of Incremental SAT Solvers [BalyoBierelserSinz'16]



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 Supports interactions during the solve () calls



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 - Notify all changes to the trail



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 - 4. Overrule found solutions
 - 5. Explain relevant propagations



Related Work

- Clingo [GebserKaminskiKaufmannOstrowskiSchaubWanko'16]
 - A state-of-the-art ASP solver
 - Supports *theory propagators*
- IntelSAT [Nadel'22]
 - Add clauses during search while maintaining propagation levels
- CP solvers [GentMiguelMoore'10]
 - □ Lazy explanation, lazy clause generation
- SAT solvers of SMT solvers [NieuwenhuisOliverasTinelli'06]
 - SAT worker interface [CimattiGriggioSchaafsmaSebastiani'13]
 - User propagators of z3 [BjørnerEisenhoferKovács'22]

Propagators are widely used, but there is no standard about how to do it.

IPASIR-UP Experiments

Extended CaDiCaL with IPASIR-UP

- A state-of-the-art incremental, inprocessing, proof producing SAT solver
- □ ~800 lines of additional code (plus another ~700 for testing)

Evaluated on two representative use cases

- Combinatorial problem solving: SAT modulo Symmetries (SMS)
- Satisfiability modulo Theories: cvc5

IPASIR-UP for SAT modulo Symmetries – Example

- Goal: Find complete set of non-isomorphic graphs over *n* vertices
 - □ Enumerate graphs with lexicographically minimal adjacency matrix
 - Would require exponential number of additional clauses
- cb_add_clause:
 - □ Enforce minimality
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			CaDiCaL+IPASIR-UP [s]			Clingo [s]	
	#vertices	#graphs	default	enum-IPASIR	no-prop	red	irred
	6	156	0.01	0.02	0.01	0.02	0.01
	7	1044	0.09	0.13	0.09	0.10	0.09
All graphs	8	12346	0.95	1.59	1.00	1.15	1.07
	9	274668	34.24	64.27	34.31	81.67	94.65
	10	12005168	50815.60	109443.72	57616.47	213959.23	196576.58

IPASIR-UP for Satisfiability Modulo Theories – cvc5

Satisfiability Modulo Theories (SMT):

- Satisfiability of a first-order formula w.r.t. some background theories
- Example theories: Arrays, bit-vectors, Arithmetic, ...
- Back-end solvers in verification, synthesis, planning, optimization, etc.
- Lazy CDCL(\mathcal{T}) approach:



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IPASIR-UP in cvc5

- cvc5: State-of-the-art SMT solver
 - □ Best ranking in several tracks of SMT'22 competition
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cvc5 with CaDiCaL through IPASIR-UP:

- Additional ~700 LOC to work with IPASIR-UP
- Encouraging performance without much tuning or optimizations
 - \square +1080 solved instances
 - $\Box ~\sim 2 imes$ faster in several logics
 - $\hfill\square$ 13 out of 19 divisions are improved
- Any IPASIR-UP supporting SAT solver can be plugged in
 - $\rightarrow~$ solid baseline to tune and improve cvc5's internals for the IPASIR-UP interface

Summary

Generic interface to inspect and influence CDCL search

- $\hfill\square$ Simple & Flexible \rightarrow relatively easy to implement
- □ Sufficient to simplify several use cases
- Implemented in a complex, modern SAT solver
 - □ Allows inprocessing of non-changing parts
- Evaluated in representative use cases (SMS, SMT)
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Future Work

- Consider further extensions
 - □ Propagate assumptions, guide backtrack, query/change variable scores, ...
- Proofs
 - □ Incremental (inprocessing) proofs
 - External proofs of external clauses
- More inprocessing

Thank you for your attention!

IPASIR Interface

// Get solver name and version **const char*** ipasir_signature(); // Initialize a solver instance and return a pointer to it **void*** ipasir_init(): // Destroy the solver instance void ipasir_release(void* solver); // Set a callback function for aborting solving void ipasir_set_terminate(void* solver, void* state, int (*terminate)(void* state)); // Add a literal or finalize clause **void** ipasir_add(**void*** solver, **int** lit_or_zero); // Assume a literal for the next solve call **void** ipasir_assume(**void*** solver, **int** lit); // Solve the formula int ipasir_solve(void* solver); // Retrieve a variables truth value (SAT case) int ipasir_val(void* solver, int lit); // Check for a failed assumption (UNSAT case) int ipasir_failed (void * solver, int lit);

Functions to Manage and Configure

```
1 // VALID = UNKNOWN | SATISFIED |
                                    UNSATISFIED
2 11
3 // require (VALID) -> ensure (VALID)
4 11
5 void connect_external propagator (ExternalPropagator * propagator);
7 // require (VALID) -> ensure (VALID)
8 11
void disconnect external propagator ():
11 // require (VALID OR SOLVING) /\ CLEAN(var) -> ensure (VALID OR SOLVING)
12 11
13 void add_observed_var (int var);
15 // require (VALID) -> ensure (VALID)
16 11
17 void remove observed var (int var):
18
19 // require (VALID OR SOLVING) -> ensure (VALID OR SOLVING)
20 11
21 bool is decision (int observed var):
23 // require (VALID OR SOLVING) -> ensure (VALID OR SOLVING)
24 11
25 void phase (int lit);
26
27 // require (VALID_OR_SOLVING) -> ensure (VALID_OR_SOLVING)
28 11
29 void unphase (int lit):
```

Example C++ implementation

```
1 class ExternalPropagator {
2 public:
    virtual ~ExternalPropagator () { }
3
4
    virtual void notify_assignment (int lit, bool is_fixed) {}
5
    virtual void notify_new_decision_level () {}
6
    virtual void notify_backtrack (size_t new_level) {}
7
8
    virtual int cb_decide () { return 0; }
9
    virtual int cb_propagate () { return 0; }
10
    virtual int cb add reason clause lit (int propagated lit) {
        return O:
12
    7
13
    virtual bool cb check found model (const std::vector<int> & model) {
14
        return true:
15
    }
16
17
    virtual bool cb has external clause () { return false: }
18
    virtual int cb_add_external_clause_lit () { return 0; }
19
20 };
```

Complete SMS Results

			CaDi	CaDiCaL+IPASIR-UP [s]			Clingo [s]	
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	#vertices	#graphs	default	no-inpro	no-prop	red	irred	
	16	0	10.58	9.14	13.58	25.07	18.56	
	17	1	39.82	31.48	44.58	122.28	87.92	
KS candida	tes 18	0	190.16	59.37	187.29	872.98	493.17	
	19	8	1220.51	1253.96	1341.80	10542.41	3348.14	
	20	147	13647.66	16449.50	13493.86	67728.42	82871.65	

Task 2: Generate up to isomorphism all non-010-colorable graphs with a minimum degree of at least three not containing a cycle of length 4 (used for Kochen-Specker Theorem)

cvc5 Results (Divisions)

		CVC5		CVC5-IPASIRUP	
_	Division	solved	time [s]	solved	time [s]
	Arith (6,865)	6,303	173,628	6,299	176,278
	BitVec (6,045)	5,552	153,899	5,529	$161,\!482$
	Equality (12,159)	5,320	2,062,804	5,322	2,061,758
	Equality+LinearArith (53,453)	45,902	2,288,230	45,906	2,288,352
	Equality+MachineArith (6,071)	983	1,533,646	987	1,532,782
	Equality+NonLinearArith (21,104)	13,314	2,419,535	13,053	2,486,588
	FPArith (3,965)	3,145	$268,\!628$	3,155	266,245
	QF_Bitvec (42,472)	40,321	984,880	40,320	985,946
	QF_Datatypes (8,403)	8,077	110,704	8,168	82,878
	QF_Equality (8,054)	8,044	9,394	8,047	7,169
	QF_Equality+Bitvec (16,585)	15,817	307,558	16,015	234,369
	QF_Equality+LinearArith (3,442)	3,388	23,041	3,381	23,465
	QF_Equality+NonLinearArith (709)	627	27,428	629	27,598
	QF_FPArith (76,238)	76,054	94,487	76,081	76,700
	QF_LinearIntArith (16,387)	11,670	1,575,635	12,004	1,512,696
	QF_LinearRealArith (2,008)	1,721	130,408	1,766	113,919
	QF_NonLinearIntArith (25,361)	13,037	4,094,712	$13,\!682$	3,840,933
	QF_NonLinearRealArith (12,134)	11,166	333,933	11,238	316,728
_	QF_Strings (69,908)	69,357	203,677	69,296	230,918
	Total (391,363)	339,798	$16,\!796,\!234$	340,878	$16,\!426,\!813$

cvc5 Results (Logics)

